

## Tutorial 5 for COMP 526 – Applied Algorithmics, Spring 2022

### Problem 1 (Fibonacci language and failure function)

The sequence of Fibonacci words  $(w_i)_{i \in \mathbb{N}_0}$  is defined recursively:

$$\begin{aligned}w_0 &= \mathbf{a} \\w_1 &= \mathbf{b} \\w_n &= w_{n-1} \cdot w_{n-2} \quad (n \geq 2)\end{aligned}$$

Unfolding the recursion yields  $w_2 = \mathbf{ba}$ ,  $w_3 = \mathbf{bab}$ ,  $w_4 = \mathbf{babba}$ , and so on.

(Note that the lengths  $|w_0|, |w_1|, |w_2|, \dots$  are *Fibonacci numbers*  $\square$ , hence the name. More precisely, we have  $|w_n| = F_{n+1}$ , with the Fibonacci numbers defined as  $F_0 = 0$ ,  $F_1 = 1$ , and  $F_n = F_{n-1} + F_{n-2}$ , for  $n \geq 2$ .)

- Construct the transition function  $\delta$  of the string-matching automaton for  $w_6$  and draw the string-matching automaton.
- Construct the prefix function  $F$  and draw the KMP automaton with failure links for  $w_6$ .

### Problem 2 (How KMP uses itself)

Recall the example  $T = \mathbf{abababaabab}$  and  $P = \mathbf{ababaca}$  used in the lecture to illustrate the KMP failure-link automaton.

- Consider the string  $S = S[0..m+n] = P\$T$  over the extended alphabet  $\Sigma' = \Sigma \cup \{\$\}$  and construct the failure-links array  $fail[0..n+m]$ .
- Compare the result with the sequence of states from simulating the failure-link automaton for  $P$  on  $T$ ; what do you observe?
- Bonus:** Can you compute the values  $fail[0..n+m]$  using only  $\Theta(P)$  extra space? Here, it is enough to have the values available at some time during the computation; we (obviously) cannot store all of them explicitly in the allowed space.