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Tutorial 8 for COMP 526 – Applied Algorithmics, Spring 2022

Problem 1 (Move-to-front transform)

Let T = T[0..9) = ABBACBAAA be an input text over alphabet $\Sigma = \{A, B, C\}$. Apply the move-to-front transform to this input with initial queue content Q = [A, B, C] and trace the content of Q throughout the execution.

Problem 2 (Lempel-Ziv-Welch compression)

Given word w = ASNXASNASNA over the ASCII character set (relevant parts of ASCII are provided on the right).

Construct, step by step, the Lempel-Ziv-Welch (LZW) factorization of w (i.e., the phrases encoded by one codeword) and provide the compressed representation of w; it suffices to show the encoded text C using integer numbers (no need for binary encodings).

Code	Character
65	A
78	N
83	S
88	Х

Problem 3 (No Free Lunch)

Prove the following *no-free-lunch* theorems for lossless compression.

1. Weak version: For every compression algorithm A and $n \in \mathbb{N}$ there is an input $w \in \Sigma^n$ for which $|A(w)| \ge |w|$, i.e. the "compression" result is no shorter than the input.

Hint: Try a proof by contradiction. There are different ways to prove this.

2. Strong version: For every compression algorithm A and $n \in \mathbb{N}$ it holds that

$$|\{w \in \Sigma^{\leq n} : |A(w)| < |w|\}| < \frac{1}{2} \cdot |\Sigma^{\leq n}|$$

In words, less than half of all inputs of length at most n can be compressed below their original size.

Hint: Start by determining $|\Sigma^{\leq n}|$.

The theorems hold for every non-unary alphabet, but you can restrict yourself to the binary case, i.e., $\Sigma = \{0, 1\}$.

We denote by Σ^* the set of all (finite) strings over alphabet Σ and by $\Sigma^{\leq n}$ the set of all strings with size $\leq n$. As domain of (all) compression algorithms, we consider the set of (all) *injective* functions in $\Sigma^* \to \Sigma^*$, i.e., functions that map any input string to some output string (encoding), where no two strings map to the same output.